

What Do Matching Models Predict?

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Testing the theory: for some statistic $S(X)$, for some F ,

$$F^S \notin (F_S^\theta).$$

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Testing: what features of the data could reject the model?

The NTU model

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Strict preferences P (everything can be rationalized under indifference)

may be represented by $U_m(w)$ for men, $V_w(m)$ for women
 0 =single, utilities $U_m(0), V_w(0)$.

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$u_m \geq U_m(0), v_w \geq V_w(0)$

and if $U_m(w) > u_m$ then $V_w(m) < v_w$ and vice-versa.

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$$u_m = \max_w \{U_m(w) | V_w(m) \geq v_w\}$$

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$\{m \mid V_w(m) \geq v_w\}$ is the *acceptance set* of woman w .

Transferable Utility

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$$u_m \geq U_m(0), v_w \geq V_w(0)$$

and:

$$\tilde{\Phi}(m, w) \equiv U_m(w) + V_w(m) \leq u_m + v_w.$$

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All men belong to the acceptance set of woman w .

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NTU: rationalizes anything *and usually much more* if (m, w) are matched, then put w at the top of P_m and vice versa.

TU: rationalizes anything as *unique equilibrium* with $\tilde{\Phi}$ in normal cone of convex polytope at the observed matching.

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(TU: theory almost-rejected since stable matching is generically unique.)

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Therefore with ≥ 3 agents on each side not every set of feasible
matchings is rationalizable.

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if H is rationalizable then it can be rationalized in **a lot** of ways.

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An *aggregate matching* is a feasible matrix $n(x, y)$ of numbers of matches per type.

E.g: $x, y = 1, 2, 3$,

$$\begin{pmatrix} 11 & 0 & 10 \\ 0 & 22 & 41 \\ 13 & 91 & 0 \end{pmatrix}$$

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More interesting:

n is rationalizable in TU iff it is rationalizable in NTU as the men-preferred or the woman-preferred matching (as with the Gale-Shapley deferred acceptance mechanism).

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But connected cycles cannot be flows.

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TU matchings cannot be a cycle for the same reason.

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then each hospital's acceptance set is a quality threshold.

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Does *not* exclude matching over unobservables; but restricts its form.

Identifying Joint Surplus

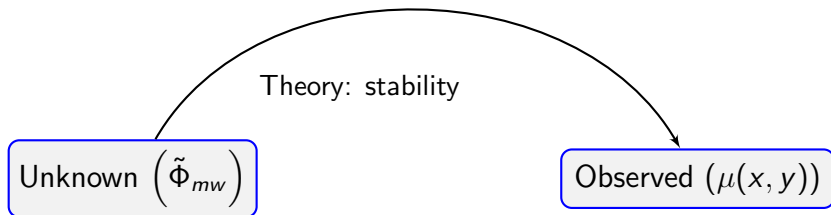
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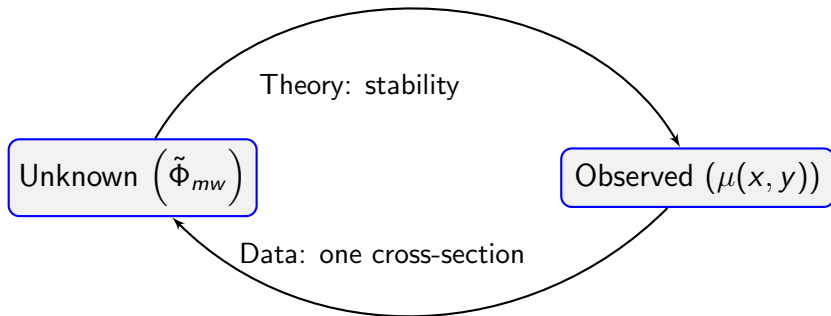
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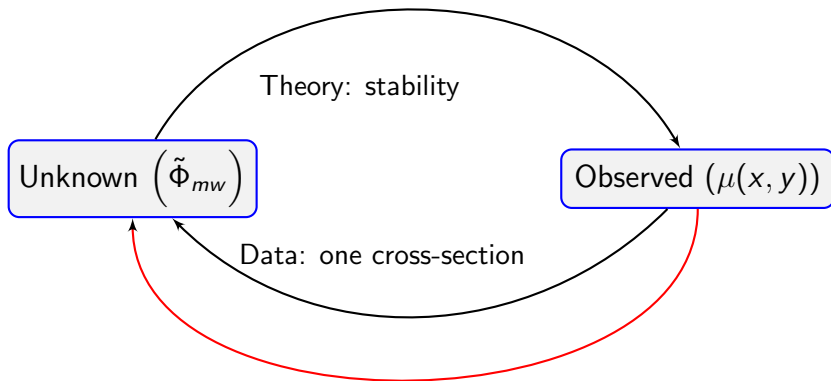
Identifying Joint Surplus



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Restrictions: separability, distributional assumptions

Consequence

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Then (Chiappori-Salanié-Weiss 2012)

Theorem

Under (S), there exists $U(x, y)$ and $V(x, y)$ such that $U(x, y) + V(x, y) = \Phi(x, y)$ and for any match $(m \in x, w \in y)$

$$u_m = U(x, y) + \varepsilon_m(y) \text{ and } v_w = V(x, y) + \eta_w(x).$$

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Proof:

$$\begin{aligned} v(w) &= \max_x \max_{m \in X} (\Phi(x, y) + \varepsilon_m(y) + \eta_w(x) - u(m)) \\ &= \max_x \left(\Phi(x, y) + \eta_w(x) - \min_{m \in X} (u(m) - \varepsilon_m(y)) \right) \\ &\equiv \max_x (\Phi(x, y) + \eta_w(x) - U(x, y)) \\ &\equiv \max_x (V(x, y) + \eta_w(x)). \end{aligned}$$

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and for each $m \in x$,
the vector $(\varepsilon_m(y))_y$ has a multidimensional cdf \mathbf{P}_x
Also assume “large markets”: many $m \in x$, many $w \in y$.

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Then **expected utilities** for a woman of type y , given all the $V(x, y)$, are

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Still many unknown quantities. . . the $U(x, y)$'s and $V(x, y)$'s.

Inverting E max for Given y

Unknown $V_{.y}$

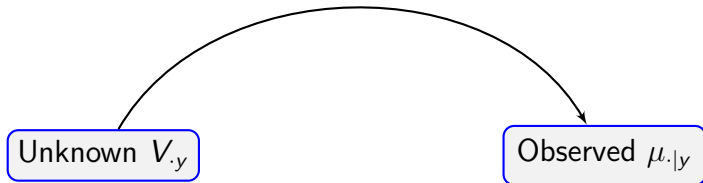
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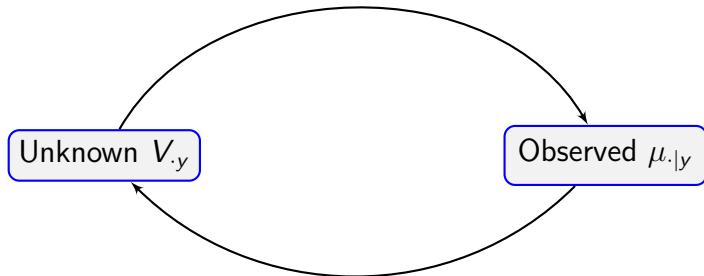
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$$\text{FOC of } H_y^*(\mu_{\cdot | y}) = \max_V (\mu \cdot V - H_y(V_{\cdot y}))$$

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define the Legendre-Fenchel (convex dual) transform:

$$f^*(y) = \max_{x \in C} (xy - f(x))$$

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=a “convex inversion formula”.

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The Legendre-Fenchel transform of H_y is

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can be estimated from the data.

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$$\Phi(x, y) = U(x, y) + V(x, y) = \frac{\partial G_x^*}{\partial \mu_{y|x}} + \frac{\partial H_y^*}{\partial \mu_{x|y}}$$

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we can even recover the full distribution of $v_w | w \in y$.

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the model is **not** testable; we have too many degrees of freedom with the error terms.

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But each of them has additional specific implications so that we can test between them.

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We need to restrict heterogeneity on pre-transfer utilities; we assume separability again,

$$U_m(w) = a(x, y) + \varepsilon_m^a(y) + \eta_w^a(x)$$

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and

$$T(x, y) = a(x, y) - \frac{\partial G_x^*}{\partial \mu_{y|x}} = \frac{\partial H_y^*}{\partial \mu_{x|y}} - b(x, y).$$

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and we can identify the distributions of ε^b and η^a if $\varepsilon^a \equiv 0$ and $\eta^b \equiv 0$,
 that is if

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Reverse question: when can we infer complementarities in surplus?

Hard if we do not know the distributions of unobserved heterogeneity;

may be possible with observed transfers in Case 2.